

NTIN071 A&G: TUTORIAL 4 – CLOSURE UNDER STRING OPERATIONS

Teaching goals: The student is able to

- formally describe a construction of an automaton based on other automata
- decide whether regular languages are closed under various string operations, including more complex ones, and prove or disprove it

IN-CLASS PROBLEMS

Problem 1 (Closure under string operations). Given a DFA A , construct an automaton recognizing the given language. (Give a formal description of the automaton.)

(a) $L(A).L(A)$	<table style="margin: auto; border-collapse: collapse;"> <tr><td style="padding: 0 5px;">→ 0</td><td style="padding: 0 5px;">a</td><td style="padding: 0 5px;">b</td></tr> <tr><td style="padding: 0 5px;">* 1</td><td style="padding: 0 5px;">1</td><td style="padding: 0 5px;">2</td></tr> <tr><td style="padding: 0 5px;">2</td><td style="padding: 0 5px;">3</td><td style="padding: 0 5px;">0</td></tr> <tr><td style="padding: 0 5px;">3</td><td style="padding: 0 5px;">4</td><td style="padding: 0 5px;">5</td></tr> <tr><td style="padding: 0 5px;">4</td><td style="padding: 0 5px;">0</td><td style="padding: 0 5px;">2</td></tr> <tr><td style="padding: 0 5px;">5</td><td style="padding: 0 5px;">2</td><td style="padding: 0 5px;">5</td></tr> <tr><td style="padding: 0 5px;">*</td><td style="padding: 0 5px;">6</td><td style="padding: 0 5px;">1</td></tr> <tr><td style="padding: 0 5px;">5</td><td style="padding: 0 5px;">5</td><td style="padding: 0 5px;">1</td></tr> <tr><td style="padding: 0 5px;">*</td><td style="padding: 0 5px;">4</td><td style="padding: 0 5px;">2</td></tr> </table>	→ 0	a	b	* 1	1	2	2	3	0	3	4	5	4	0	2	5	2	5	*	6	1	5	5	1	*	4	2	<table style="margin: auto; border-collapse: collapse;"> <tr><td style="padding: 0 5px;">→ 0</td><td style="padding: 0 5px;">a</td><td style="padding: 0 5px;">b</td></tr> <tr><td style="padding: 0 5px;">* 1</td><td style="padding: 0 5px;">0</td><td style="padding: 0 5px;">5</td></tr> <tr><td style="padding: 0 5px;">2</td><td style="padding: 0 5px;">1</td><td style="padding: 0 5px;">3</td></tr> <tr><td style="padding: 0 5px;">3</td><td style="padding: 0 5px;">2</td><td style="padding: 0 5px;">5</td></tr> <tr><td style="padding: 0 5px;">*</td><td style="padding: 0 5px;">3</td><td style="padding: 0 5px;">2</td></tr> <tr><td style="padding: 0 5px;">*</td><td style="padding: 0 5px;">4</td><td style="padding: 0 5px;">6</td></tr> <tr><td style="padding: 0 5px;">5</td><td style="padding: 0 5px;">5</td><td style="padding: 0 5px;">1</td></tr> <tr><td style="padding: 0 5px;">*</td><td style="padding: 0 5px;">6</td><td style="padding: 0 5px;">4</td></tr> <tr><td style="padding: 0 5px;">*</td><td style="padding: 0 5px;">4</td><td style="padding: 0 5px;">2</td></tr> </table>	→ 0	a	b	* 1	0	5	2	1	3	3	2	5	*	3	2	*	4	6	5	5	1	*	6	4	*	4	2
→ 0	a	b																																																						
* 1	1	2																																																						
2	3	0																																																						
3	4	5																																																						
4	0	2																																																						
5	2	5																																																						
*	6	1																																																						
5	5	1																																																						
*	4	2																																																						
→ 0	a	b																																																						
* 1	0	5																																																						
2	1	3																																																						
3	2	5																																																						
*	3	2																																																						
*	4	6																																																						
5	5	1																																																						
*	6	4																																																						
*	4	2																																																						
(b) $L(A)^+$																																																								
(c) $L(A)^*$																																																								
(d) $L(A)^R$																																																								

Problem 2 (Delete). Let L be a regular language over the alphabet $\Sigma = \{a, b\}$. Describe the following languages in set notation. Decide if they are (necessarily) also regular, prove or disprove. The language of all words obtained from words of the language L by...

- ... deleting all occurrences of the letter a .
- ... deleting the initial letter and writing this letter at the end of the word.
- ... deleting the longest contiguous sequence of a 's from the beginning of the word.

EXTRA PRACTICE AND THINKING

Problem 3 (Prefixes). Are regular languages closed under the following operations? Prove or disprove. (In the following, L is a regular language over an alphabet Σ .)

- $\text{init}(L) = \{w \in \Sigma^* \mid \text{there is } u \in \Sigma^* \text{ such that } wu \in L\}$
- $\text{min}(L) = \{w \in L \mid \text{there is no } u \in L, v \in \Sigma^+ \text{ such that } w = uv\}$
- $\text{max}(L) = \{w \in L \mid \text{there is no } u \in \Sigma^+ \text{ such that } wu \in L\}$

Problem 4 (Shift). Given a regular language L over an alphabet Σ , define the language L' as follows. Is the language L' necessarily regular?

$$L' = \{uv \mid u, v \in \Sigma^*, vu \in L\}$$

Problem 5 (Cut). Consider two regular languages L and M over an alphabet Σ , and define the language K as follows. Is the language K necessarily regular?

$$K = \{uw \mid u, w \in \Sigma^*, (\exists v \in M) uvw \in L\}$$

Problem 6 (Iterations of unary languages). Show that for any language L over the alphabet $\Sigma = \{a\}$, the language L^* is regular.