

NTIN071 A&G: TUTORIAL 6 – FORMAL GRAMMARS, REGULAR AND CONTEXT-FREE GRAMMARS

Teaching goals: The student is able to

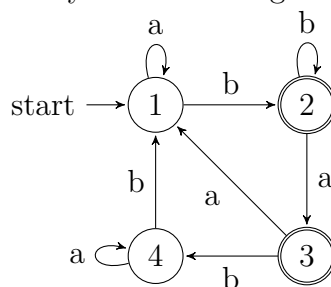
- explain the formal definition of a grammar, and the language it generates,
- give definitions and examples grammars of all types in the Chomsky hierarchy,
- describe a language generated by a given context-free grammar,
- construct a grammar for a language given in set notation,
- convert a finite automaton to a right-linear grammar,
- convert a right-linear grammar to a finite automaton,
- design algorithms to test basic properties of context-free grammars.

IN-CLASS PROBLEMS

Problem 1 (Constructing grammars). Design grammars (of the highest possible type) which generate the following languages ($\Sigma = \{a, b\}$ unless specified otherwise):

- | | |
|---|--|
| (a) $L = \{w \in \Sigma^* \mid w _b \text{ is even}\}$ | (d) $L = \{a^i b^j \mid 0 \leq i \leq j \leq 2i\}$ |
| (b) $L = \{ww^R \mid w \in \Sigma^*\}$ | (e) $L = \{uabbav \mid u, v \in \Sigma^* \text{ and } u = v \}$ |
| (c) $L = \{a^i b^j c^k \mid i = j \text{ or } j = k\}$ | (f) $L = \{w \in \Sigma^* \mid -1 \leq w _a - w _b \leq 1\}$ |

Problem 2 (FA to grammar). For the following automaton, find an equivalent grammar. Which class of the Chomsky hierarchy does it belong to?



Problem 3 (Regular grammar to FA). Convert the following right-linear grammar to a finite automaton: $G = (\{S, A, B, C\}, \{a, b\}, \mathcal{P}, S)$ where \mathcal{P} consists of the following:

$$\begin{aligned}
 S &\rightarrow abS \mid babA \mid \epsilon \\
 A &\rightarrow abA \mid aB \mid bC \\
 B &\rightarrow abS \mid B \mid bC \mid \epsilon \\
 C &\rightarrow aab \mid A \mid aA \mid \epsilon
 \end{aligned}$$

Problem 4 (Testing properties of context-free languages). Design an (efficient) algorithm which decides whether a given context-free grammar satisfies the given property:

- | | | |
|---------------------------|-------------------------|---------------------------------|
| (a) $L(G) \neq \emptyset$ | (b) $\epsilon \in L(G)$ | (c) $L(G)$ is a finite language |
|---------------------------|-------------------------|---------------------------------|

EXTRA PRACTICE AND THINKING

Problem 5 (Constructing grammars). Design grammars (of the highest possible type) which generate the following languages ($\Sigma = \{a, b\}$ unless specified otherwise):

- (a) $L = \Sigma^*$
- (b) $L = \{a^{2i}b^j \mid i \leq j\}$
- (c) $L = \{w \in \Sigma^* \mid |w|_a = 2|w|_b\}$
- (d) $L = \{uabbav \mid u, v \in \Sigma^* \text{ and } |u| \neq |v|\}$
- (e) $L = \{w\#s^R \mid w, s \in \Sigma^* \text{ and } s \text{ is a subword of } w\}$

Problem 6 (Small grammars generating large (finite) languages). Find a sequence of context-free grammars G_1, G_2, G_3, \dots (over a given alphabet Σ) such that G_n generates exactly all words of length $\leq 2^n$ (and no other words), and the size of G_n (for simplicity, say the number of symbols in bodies of production rules) is in $O(n)$.

Problem 7 (Context-sensitive Grammar?). Let $G = (\{S, A, B, C\}, \{a, b, c\}, \mathcal{P}, S)$, where:

$$\mathcal{P} = \{S \rightarrow aSBC \mid aBC, B \rightarrow BBC, C \rightarrow CC, CB \rightarrow BC, \\ aB \rightarrow ab, bB \rightarrow bb, bC \rightarrow bc, cC \rightarrow cc\}$$

What language does it generate? Is the grammar G context-sensitive? If not, find an equivalent context-sensitive grammar.