

NTIN071 A&G: TUTORIAL 7 – PUMPING LEMMA FOR CONTEXT-FREE LANGUAGES,
CHOMSKY NORMAL FORM

Teaching goals: The student is able to

- give a formal statement of the Pumping lemma for context-free languages
- apply the Pumping lemma to prove that a given language is not context-free
- give the formal definition of Chomsky Normal Form and related notions
- explain the proof of the Pumping lemma for context-free languages
- convert a given context-free grammar to ChNF
- explain the CYK algorithm, apply to a given word and context-free grammar

IN-CLASS PROBLEMS

Problem 1 (Pumping lemma: statement and proof). (a) Formulate the Pumping Lemma for context-free languages (without consulting your notes).

(b) Compare the statement to the version for regular languages.

(c) Explain the idea behind its proof.

(d) Demonstrate pumping on the language $L = \{ww^R \mid w \in \{a, b\}^*\}$.

Problem 2 (Pumping lemma: application). Decide if the following languages are context-free. Prove that your answer is correct.

(a) $L = \{0^i 1^i \mid i \geq 0\}$

(e) $L = \{ww^R \mid w \in \{0, 1\}^*\}$

(b) $L = \{0^i 1^j 0^i \mid 0 \leq j \leq i\}$

(f) $L = \{1^{n^2} \mid n \geq 0\}$

(c) $L = \{0^i 1^i 2^i \mid i \geq 0\}$

(g) $L = \{1^p \mid p \text{ is a prime}\}$

(d) $L = \{0^{2^i} 1^{3^i} 0^i \mid i \geq 0\}$

(h) $L = \{0^i 1^j \mid 0 \leq i \leq j^2\}$

Problem 3 (Convert to ChNF). Convert the following context-free grammars to Chomsky normal form:

(a) $G_1 = (\{S, A, B\}, \{0, 1\}, S, \mathcal{P})$, where

(b) $G_2 = (\{S, A, B\}, \{0, 1\}, S, \mathcal{P})$, where

$$\mathcal{P} = \{S \rightarrow 0AB,$$

$$\mathcal{P} = \{S \rightarrow 0A10B10,$$

$$A \rightarrow 0A0 \mid 11,$$

$$A \rightarrow 1A0 \mid \epsilon,$$

$$B \rightarrow 0\}$$

$$B \rightarrow 1B00 \mid \epsilon\}$$

Problem 4 (The CYK algorithm). Using the CYK algorithm determine if $w \in L(G)$.

(a) $w = 001100$, $G = G_1$ is the grammar from Problem 3(a)

(b) $w = 110011$, $G = G_1$ is the grammar from Problem 3(a)

EXTRA PRACTICE AND THINKING

Problem 5 (Pumping lemma: application). Decide if the following languages are context-free. Prove that your answer is correct.

- (a) $L = \{0^i 1^j 0^i \mid i, j \geq 0\}$ (d) $L = \{ww \mid w \in \{0, 1\}^*\}$
 (b) $L = \{0^i 1^j 0^i \mid 0 \leq i \leq j\}$ (e) $L = \{ww^R \mid w \in \{0, 1\}^*, |w|_0 = |w|_1\}$
 (c) $L = \{0^i 1^j 2^k \mid 0 \leq i \leq j \leq k\}$ (f) $L = \{1^{n^2+n+1} \mid n \geq 0\}$

Problem 6 (Pumping and right-linear grammars). Give an alternative proof of the Pumping lemma for regular languages that is based on derivations from a right-linear grammar.

Problem 7 (Pumping linear languages). Recall that a grammar is *linear*, if it only contains production rules of the form $A \rightarrow uBw$ and $A \rightarrow w$, where $A, B \in V$ and $u, w \in T^*$.

- (a) Formulate a Pumping lemma for linear languages.
 (b) Proof the statement using derivations from a (reduced) linear grammar.
 (c) How does the pumping constant n from the lemma relate to a linear grammar for L ?
 (d) Show that the language $L = \{w \in \{0, 1\}^* \mid |w|_0 = |w|_1\}$ is not linear.
 (e) Where does L lie within the Chomsky hierarchy?

Problem 8 (About the conversion to ChNF). Recall the process of converting a context-free grammar to Chomsky Normal Form. Then answer the following questions. Justify.

- (a) Find an example of a grammar in which there is a generating variable only reachable via nongenerating variables.
 (b) When reducing a grammar, which variables do we need to remove first: nongenerating or unreachable?
 (c) Is it possible for a reachable generating variable to become nongenerating after the removal of unreachable variables?
 (d) When we want to break up a production rule with long body, what is the minimal number of Chomsky Normal Form rules we need to create?

Problem 9 (Convert to ChNF). Convert the following to Chomsky normal form:

- (a) $G = (\{S, A, B\}, \{0, 1\}, S, \mathcal{P})$ (b) $G = (\{S, E, F\}, \{(\cdot), *, +, \cdot, 1\}, S, \mathcal{P})$
 $\mathcal{P} = \{S \rightarrow A \mid 0SA \mid \epsilon,$ $\mathcal{P} = \{S \rightarrow (E),$
 $A \rightarrow 1A \mid 1 \mid B1,$ $E \rightarrow F + F \mid F * F,$
 $B \rightarrow 0B \mid 0 \mid \epsilon\}$ $F \rightarrow S \mid 1\}$

Problem 10 (The CYK algorithm). Using the CYK algorithm determine if $w \in L(G)$.

- (a) $w = 01010010$, $G = G_2$ is the grammar from Problem 3(b)
 (b) $w = 01010011$, $G = G_2$ is the grammar from Problem 3(b)